

Reversible ALU Circuit Realization Using Verilog HDL

Vishal A. Wankhede, Ramprasad M. Gawande

Department of E&TC, SNJB's CoE, Chandwad (Nasik), India

E-mail: rampgawande611@gmail.com

Abstract

Reversibility in computing implies that no info regarding the process states will ever be lost, thus, we are able to recover any earlier stage by computing backwards or un-computing the results. This is often termed as logical changeability. The advantages of logical changeability are gained solely when using physical changeability. Physical changeability may be a method that dissipates no energy to heat. Completely excellent physical changeability is much impossible. Computing systems provide off heat once voltage levels modification from positive to negative: bits from 0 to 1. Most of the energy required to create that modification is given off within the style of heat. Rather than dynamic voltages to new levels, reversible circuit parts can bit by bit move charge from one node to subsequent. This way, one will solely expect to lose a second quantity of energy on every transition. Reversible computing powerfully affects digital logic styles. Reversible logic parts square measure required to recover the state of inputs from the outputs. It will impact instruction sets and high-level programming languages similarly. Eventually, these will need to be reversible to produce best potency.

Keywords: *Digital logic designs, logical reversibility, optimal efficiency, physical reversibility*

INTRODUCTION

High-performance chips cathartic massive amounts of warmth impose sensible limitation on, however, so much will we have a tendency to improve the performance of the system. Reversible circuits that conserve info, by un-

computing bits rather than throwing them away, can before long provide the sole physically attainable thanks to keep up performance. Reversible computing also will result in improvement in energy potency. Energy potency can basically have an effect on the speed of circuits

appreciate nano-circuits and thus the speed of most computing applications. To extend the movableness of devices once more reversible computing is needed [1]. It can let circuit part sizes to cut back to atomic size limits and, therefore, devices will become a lot of moveable. Though the hardware style prices incurred in close to future could also be high, however, the facility value and performance being a lot of dominant than logic hardware value in today's computing era, the requirement of reversible computing cannot be unheeded.

A reversible computer circuit is associate in nursing n-input n-output logic device with matched mapping. This helps to work out the outputs from the inputs and additionally the inputs will be unambiguously recovered from the outputs. Additionally, within the synthesis of reversible circuits direct fan-out is not allowed as one-to-many idea is not reversible. But fan-out in reversible circuits is achieved victimisation extra gates [2, 3]. A reversible circuit ought to be designed victimisation minimum variety of reversible logic gates. From the

purpose of read of reversible circuit style, there are several parameters for crucial the quality and performance of circuits.

EXISTING SYSTEM

In existing Design Arithmetic Logic Unit is designed with the help of Conventional or irreversible Logic Gates [4, 5]. It has been known for a long time that the two-bit gates, AND & OR, and the one-bit gate NOT, are universal for classical computation, in the sense that they are sufficient to simulate any function. Gate functions in classical logic are often represented using truth tables. The AND & OR gates have two inputs and one output, while the NOT gate has one input and one output. The DM74LS181 is reference logic for proposed design, which is a 4-bit ALU and can perform all the possible 16 logic operations on two variables and a variety of arithmetic operations. ALU provides 16 arithmetic operations: add, subtract, compare, double, plus twelve other arithmetic operations. Provides 16 logic operations of two variables: EXOR, compare, AND, NAND, OR, NOR, plus ten other logic operations [6, 7].

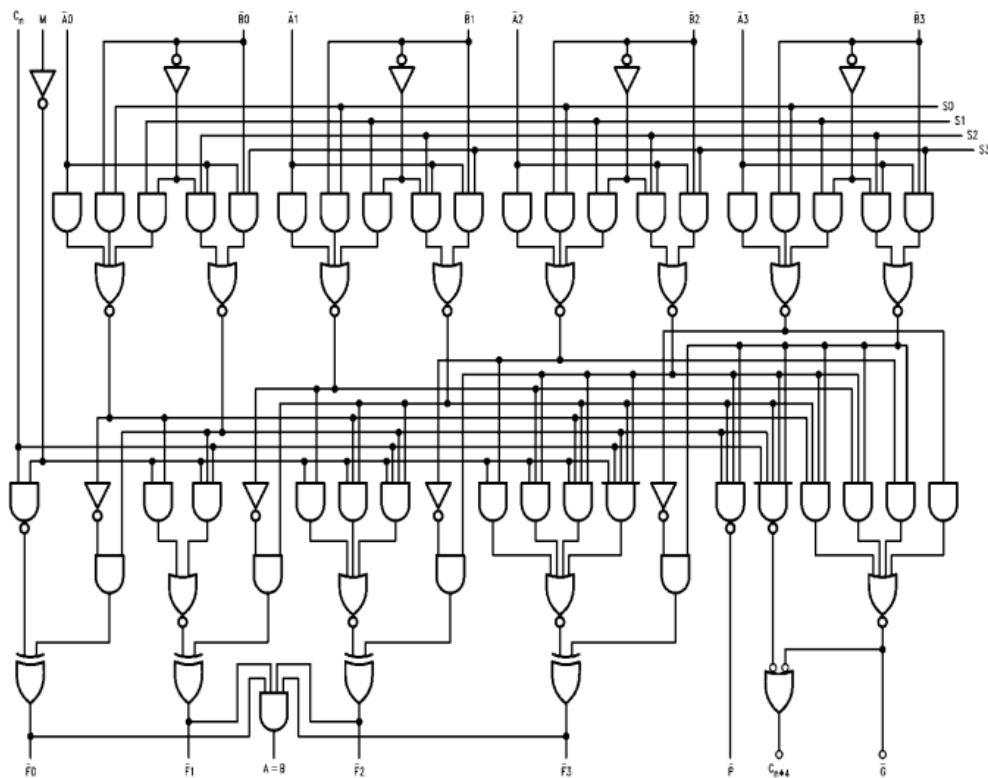


Fig. 1: Existing Conventional ALU.

Existing System Drawbacks

- More Complexity.
- Costlier Circuits.
- Delayed Operations.
- Not Power Efficient.

PROPOSED SYSTEM

A reversible gate has equal variety of input and output terminals and there is one to at least one mapping between them. Once more we are able to say, gate is reversible if we are able to verify input vector from output vector and the other way around.

Reversible gate ought to much loose little or no quantity of energy. Fan-out is not allowed in reversible circuits but fan-out is often achieved victimisation further gate. During this paper, we have to implement associate degree eight Bit Arithmetic Logic Unit with the assistance of basic reversible gate like Richard Phillips Feynman gate, Fredkin gate, NFT gate, NEW gate, Peres gate, Toffoli gate and URG gate. We will implement reversible ordered circuits based mostly ALU [8–10].

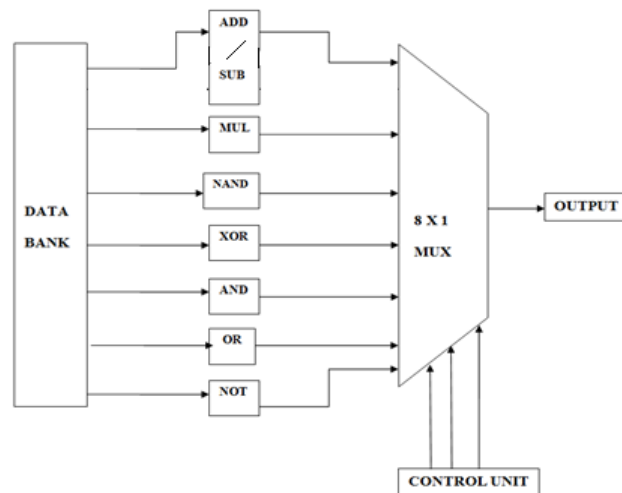


Fig. 2: Proposed Reversible ALU.

The multi-function ALU supported reversible logic gates in the main contains the reversible operate generator (FUNC) and, therefore, the reversible controlled unit (DXOR). The reversible operate generator and, therefore, the reversible controlled unit are cascaded by some n-Toffoli gates and NOT gates, and whimsical bit reversible ALU modules may be accomplished by this manner. Within the procedure of cascading the reversible operate generator and, therefore, the reversible controlled unit, we tend to

utilise the output signals to scale back the value of circuit style the maximum amount as potential.

Feynman Gate

The Feynman gate also known as a Controlled NOT (CNOT) gate is one of the popular examples of a 2 x 2 reversible gate. As shown in Figure 3, the first input in this gate is passed to the output without any change and the second output is the XOR of the first and second inputs.

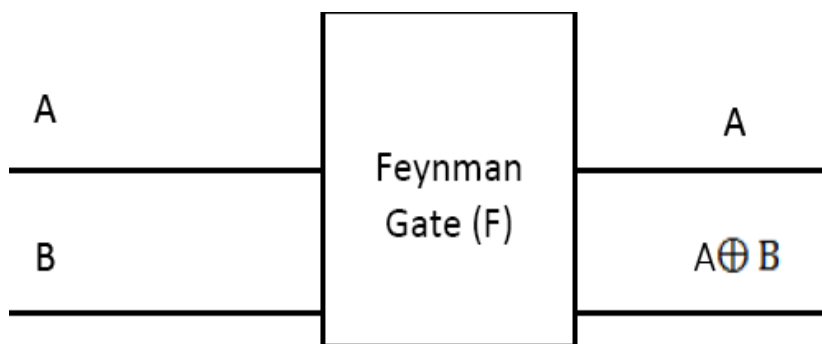


Fig. 3: Feynman Gate.

$$X = a$$

$$Y = a \text{ XOR } b$$

Toffoli Gate

The Toffoli gate, or a doubly controlled NOT gate, is a (3X3) universal reversible gate. As shown in Figure 4, the first two inputs are directly passed to the

corresponding outputs and the third output is the logical XOR of the third input with the logical AND of all the first two inputs. a, b and c are the three inputs to the gate and the corresponding Output lines are X, Y and Z respectively. The corresponding functions computed are as following:

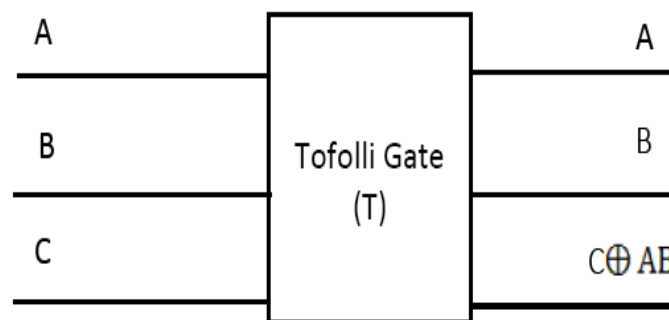


Fig. 4: Toffoli Gate.

$$X = a$$

$$Y = b$$

$$Z = c \text{ XOR } ab$$

As the traditional method of reversible gate design was too lengthy and ineffective, a new approach was adopted. The operations required to be performed by the ALU were taken- AND, XOR, NOR, and SUM. A circuit was made for each operation separately and it was then combined in such a manner so as to get the

minimum line and gate cost. A 4:1 reversible multiplexer was used to get the output at one line. The multiplexer utilizes Fredkin gates. If modified Fredkin gates are used for the multiplexer, the quantum cost can be decreased further. The AXONS ALU, shown in Figure 5, was named after the operations it performs – And, XOR, NOR, Sum. It has two select lines and operates according to the truth table.

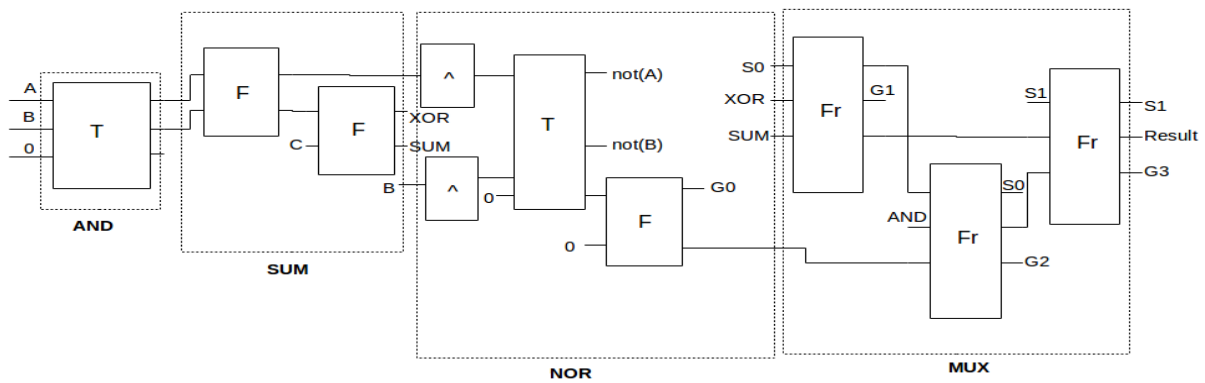


Fig. 5: Proposed ALU Design.

RESULTS

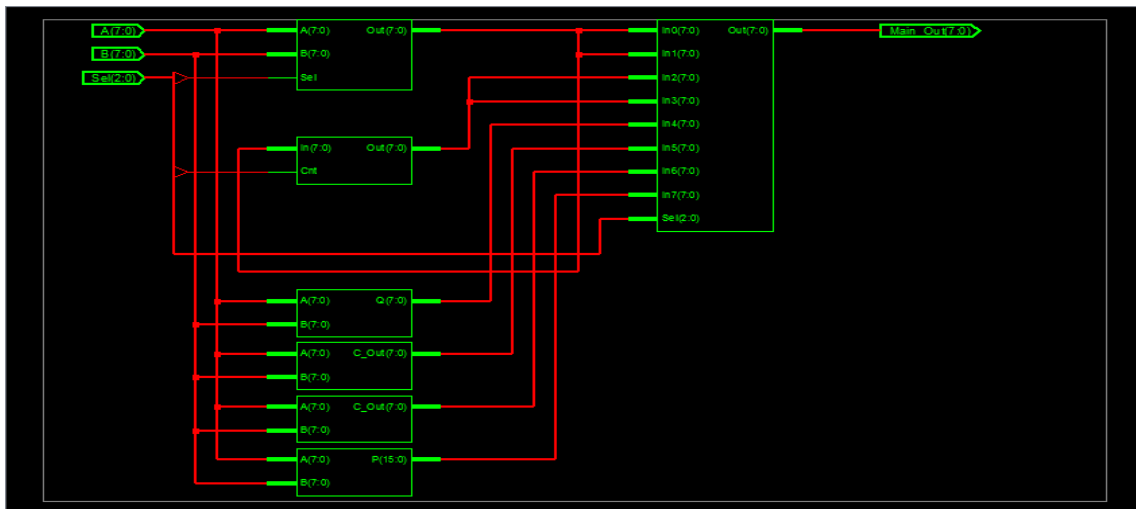


Fig. 6: ALU RTL Schematic.

Messages					
/Main_8_Bit_ALU/A	00010101	00010101			
/Main_8_Bit_ALU/B	00001010	00001010			
/Main_8_Bit_ALU/Sel	111	1011	100	101	110
/Main_8_Bit_ALU/Main_Out	11010010	00001010	00011111	00000000	00011111
/Main_8_Bit_ALU/O1	00011111	00011111			
/Main_8_Bit_ALU/O2	00000000	00000000			
/Main_8_Bit_ALU/O3	00001011	00001011	00100000	00001011	00100000
/Main_8_Bit_ALU/O4	00001010	00001010	00100001	00001010	00001010
/Main_8_Bit_ALU/O5	00011111	00011111			
/Main_8_Bit_ALU/O6	0000000011010010	0000000011010010			
/Main_8_Bit_ALU/SD0/A	00010101	00010101			
/Main_8_Bit_ALU/SD0/B	00001010	00001010			
/Main_8_Bit_ALU/SD0/Q	00011111	00011111			
/Main_8_Bit_ALU/SD0/P	00010101	00010101			

Fig. 7: ALU Output Bit for given Input.

Table 1: Comparison of Area, Delay, Power.

Method Name	Area in Number of LUT			Delay			Power
	LUT	Gate Counts	Slices	Delay	Logic Delay	Path Delay	
Conventional ALU	136	873	72	26.711ns	13.507ns	13.204ns	43 mW
Reversible ALU	120	759	65	24.620 ns	12.886ns	11.734ns	41 mW

Proposed System Advantages

- Complexity Reduced.
- Speed Increased.
- Delay Minimized.
- Reduction in Cost.

Real Time Application

- Fault Coverage for Single Missing/Additional cell.
- Testing Sequential Circuits.

CONCLUSION

This paper planned reversible successive circuits supported conservative logic that is checkable for any unifacial stuck-at faults mistreatment solely two tests vectors, all 0s and all 1s. The planned successive circuits supported conservative logic gates surmount the successive circuit enforced in classical gates in terms of testability. The successive circuits enforced mistreatment typical classic gates do not offer genetic support for testability. Hence, a traditional successive circuit

wants modification within the original electronic equipment to supply the testing capability. Additionally, because the quality of a successive circuit will increase the quantity of check vector needed to check the successive circuit additionally will increase. For example, check a posh consecutive circuit thousand of test vectors square measure needed to check all stuck-at-faults, whereas, if identical consecutive circuit is build mistreatment planned reversible consecutive building blocks it will be checked by solely two test vectors, all 0s and all 1s. Thus, the most advantage of the planned conservative reversible consecutive circuits compared to the standard consecutive circuit is that they would like of solely two check vectors to check any consecutive circuit no matter its complexity.

REFERENCES

1. J. Ren, V. K. Semenov. Progress with physically and logically reversible

- superconducting digital circuits. *IEEE Trans. Appl. Superconduct.* 2011; 21(3): 780–786p.
2. S. F. Murphy, M. Ottavi, M. Frank, et al. On the design of reversible QDCA systems. *Sandia National Laboratories, Albuquerque, NM, Tech. Rep. SAND2006-5990*; 2006.
 3. H. Thapliyal, N. Ranganathan. Reversible logic-based concurrently testable latches for molecular QCA. *IEEE Trans. Nanotechnol.* 2010; 9(1): 62–69p.
 4. P. Tougaw, C. Lent. Logical devices implemented using quantum cellular automata. *J. Appl. Phys.* 1994; 75(3): 1818–1825p.
 5. P. Tougaw, C. Lent. Dynamic behavior of quantum cellular automata. *J. Appl. Phys.* 1996; 80(8): 4722–4736p.
 6. M. B. Tahoori, J. Huang, M. Momenzadeh, et al. Testing of quantum cellular automata. *IEEE Trans. Nanotechnol.* 2004; 3(4): 432–442p.
 7. G. Swaminathan, J. Aylor, B. Johnson. Concurrent testing of VLSI circuits using conservative logic. *In Proc. Int. Conf. Comput. Design.* Cambridge, MA; 1990: 60–65p.
 8. E. Fredkin, T. Toffoli. Conservative logic. *Int. J. Theor. Phys.* 1982; 21(3–4): 219–253p.
 9. P. Kartschoke. Implementation issues in conservative logic networks. Charlottesville; 1992.
 10. G. Swaminathan. Concurrent error detection techniques using parity. Charlottesville; 1989.